Doctoral Thesis in Applied and Computational Mathematics

Numerical methods for parameterized linear systems

SIOBHÁN CORRENTY

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Doctoral Thesis in Applied and Computational Mathematics
KTH Royal Institute of Technology
Stockholm, Sweden 2024
To my family and friends
Abstract

Solving linear systems of equations is a fundamental problem in engineering. Moreover, applications involving the solution to linear systems arise in the social sciences, business, and economics. Specifically, the research conducted in this dissertation explores solutions to linear systems where the system matrix depends nonlinearly on a parameter. The parameter can be a scalar or a vector, and a change in the parameter results in a change in the solution. Such a setting arises in the study of partial differential equations and time-delay systems, and we are interested in obtaining solutions corresponding to many values of the parameter simultaneously. The methods developed in this thesis can also be used to solve parameter estimation problems. Furthermore, software has been developed and is available online.

This thesis consists of four papers and presents both algorithms and theoretical analysis. In Paper A, a linearization based on an infinite Taylor series expansion is considered. Specifically, the linearized system is a shifted parameterized system, and the parameter is a scalar. The GMRES method is used to solve the systems corresponding to many values of the parameter, and only one Krylov subspace basis matrix is required. Convergence analysis is based on solutions to a nonlinear eigenvalue problem and the magnitude of the parameter. Notably, the algorithm is carried out in a finite number of computations.

The approach in Paper B is based on a preconditioned linearized system solved using the inexact GMRES method. In this setting, the linearization incorporates all terms in an infinite Taylor series expansion, and the preconditioner is applied approximately using iterative methods. Solutions corresponding to many values of the scalar parameter are generated from one subspace, and this is done in a finite number of linear algebra operations. Theoretical analysis, based on the error in the application of the preconditioner and the magnitude of the parameter, leads to a bound on the residual.

Paper C proposes a short recurrence Krylov subspace method for solving linear systems that depend on a scalar parameter. In particular, a Chebyshev approximation is used to construct a linearization, and the linearized system is solved in a Bi-CG setting. Additionally, shift-and-invert preconditioning leads to fast convergence of the Krylov method for many different values of the parameter. An inexact variant of the method is also derived and analyzed.

In Paper D, a reduced order model is constructed from snapshots to solve parameterized linear systems. Specifically, the parameter is a vector of dimension 2, and the sampling is performed on a sparse grid using the method proposed in Paper C. A tensor decomposition is utilized to build the model. Approaches of this kind are not always successful, and it is not known a priori if a decomposition will converge on a given set of snapshots. This work offers a novel way to generate a new set of snapshots in the same parameter space, to be used if the decomposition does not converge, with little extra computation.
Keywords: Parameterized linear systems, Krylov subspace methods, preconditioning, tensor decompositions, shifted linear systems, parameterized partial differential equations, time-delay systems, transfer functions, parameter estimation problems


Tillvägagångssättet i artikel B är baserat på ett förkonditionerat linjäriserat system löst med den inexakta GMRES-metoden. I den här contexten innehåller linjäriseringen alla termer i en oändlig Taylor-serieexpansion, och förkonditioneringen appliceras på ett approximativt sätt med iterativa metoder. Lösningar som motsvarar många värden på den skala parameter genereras från ett delrum, och detta görs i ett ändligt antal linjära algebraoperationer. Teoretisk analys, baserad på felet i appliceringen av förkonditioneraren och storleken på parametern, leder till en övre begränsning på residualens storlek.


I artikel D konstrueras en reducerad ordningsmodell från sampel av modellen för att lösa parametraserade linjära system. Specifikt är parameteren en vektor med dimensionen 2, och samplingen utförs på ett glest rutnät med den metod som föreslås i artikel C. En tensorfaktorisering används för att bygga modellen. Tillvägagångssättet av denna typ är inte alltid framgångsrika, och det är inte känt på förhand om en tensorfaktorisering kommer att konvergera för en given uppsättning av sampel. Detta arbete presenterar ett nytt sätt att generera en ny uppsättning sampel i samma parameterrum till en låg extra kostnad. De nya lösningar kan användas om tensorfaktoriseringen misslyckas.
Nyckelord: Parameteriserade linjära system, Krylov-metoder, förkonditionering, tensordekomposition, skiftade linjära system, parametriserade partiella differentialekvationer, tidsfördröjningssystem, överföringsfunktioner, parameteruppskattningsproblem
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Stockholm, May 2024
Siobhán Correnty
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Paper B
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Paper C
*Preconditioned Chebyshev BiCG method for parameterized linear systems*

Paper D
*Chebyshev HOPGD with sparse grid sampling for parameterized linear systems*
## Notations and Abbreviations

<table>
<thead>
<tr>
<th>Notation</th>
<th>Description</th>
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<tbody>
<tr>
<td>( \mathbb{R} )</td>
<td>Set of real numbers</td>
</tr>
<tr>
<td>( \mathbb{R}^n )</td>
<td>Space of real column vectors of dimension ( n )</td>
</tr>
<tr>
<td>( \mathbb{R}^{n \times m} )</td>
<td>Space of real matrices of dimension ( n \times m )</td>
</tr>
<tr>
<td>( \mathbb{R}^{n \times m_1 \times m_2} )</td>
<td>Space of real tensors of dimension ( n \times m_1 \times m_2 )</td>
</tr>
<tr>
<td>( \mathbb{R}^+ )</td>
<td>Set of positive real numbers</td>
</tr>
<tr>
<td>( \mathbb{C} )</td>
<td>Set of complex numbers</td>
</tr>
<tr>
<td>( \mathbb{Z} )</td>
<td>Set of integers</td>
</tr>
<tr>
<td>( \mathbb{N} )</td>
<td>Set of natural numbers 1, 2, \ldots</td>
</tr>
<tr>
<td>(</td>
<td>\alpha</td>
</tr>
<tr>
<td>( a_{i,j} )</td>
<td>Entry ((i,j)) of a matrix ( A )</td>
</tr>
<tr>
<td>( a_j )</td>
<td>( j )th column of a matrix ( A )</td>
</tr>
<tr>
<td>( A(1:i,1:j) )</td>
<td>Rows 1 to ( i ), columns 1 to ( j ) of a matrix ( A )</td>
</tr>
<tr>
<td>( A^T )</td>
<td>Transpose of a matrix ( A )</td>
</tr>
<tr>
<td>( A^* )</td>
<td>Hermitian transpose of a matrix ( A )</td>
</tr>
<tr>
<td>( K_k(A,b) )</td>
<td>Krylov subspace of dimension ( k ) generated from the matrix ( A \in \mathbb{C}^{n \times n} ) and the vector ( b \in \mathbb{C}^n )</td>
</tr>
<tr>
<td>( \dim(K_k(A,b)) )</td>
<td>Dimension of ( K_k(A,b) )</td>
</tr>
<tr>
<td>( \text{tr}(A) )</td>
<td>Trace of matrix ( A )</td>
</tr>
<tr>
<td>( \text{vec}(A) )</td>
<td>Vectorization of matrix ( A )</td>
</tr>
<tr>
<td>( \sigma_i(A) )</td>
<td>( i )th largest singular value of ( A )</td>
</tr>
<tr>
<td>( \sigma_{\text{min}}(A), \sigma_{\text{max}}(A) )</td>
<td>Smallest and largest singular value of ( A )</td>
</tr>
<tr>
<td>( \lambda(A) )</td>
<td>Set of eigenvalues of ( A )</td>
</tr>
<tr>
<td>( \lambda_i(A) )</td>
<td>( i )th largest eigenvalue of ( A )</td>
</tr>
<tr>
<td>( \text{diag}(D_1,\ldots,D_n) )</td>
<td>Block matrix with ( D_i ) on the diagonal</td>
</tr>
<tr>
<td>( \partial \Omega )</td>
<td>Boundary of a set ( \Omega )</td>
</tr>
<tr>
<td>( I_n )</td>
<td>Identity matrix of dimension ( n \times n )</td>
</tr>
</tbody>
</table>
Notations and Abbreviations

\( 0_{n \times m} \) Zero matrix of dimension \( n \times m \)
\( \otimes \) Kronecker product
\( \|x\|, \|x\|_p \) 2-norm and \( p \)-norm of a vector \( x \)
\( \|A\|, \|A\|_p \) 2-norm and \( p \)-norm of a matrix \( A \)
\( e_i \in \mathbb{R}^n \) \( i \)th column of \( I_n \)
\( \mathcal{N}(\mu, \sigma^2) \) Gaussian distribution with mean \( \mu \) and variance \( \sigma^2 \)
\( \mathbb{E}(X) \) Expectation of a random variable \( X \)
\( \kappa(A), \kappa_p(A) \) 2-norm and \( p \)-norm condition number of \( A \)
\( \text{span}\{v_1, \ldots, v_n\} \) Span of the set of vectors \( v_1, \ldots, v_n \)
\( (u, v)_\Omega \) Integral of the dot product of \( u, v \in \mathbb{C}^n \) over \( \Omega \)
\( \inf(S), \sup(S) \) Infimum and supremum of a set \( S \)
\( (K_N - \mu M_N)v(\mu) = c \) Linearization of \( A(\mu)x(\mu) = b \)

Abbreviations

ADI Alternating directions implicit
AGMG Aggregation-based algebraic multigrid
Bi-CG Bi-conjugate gradient
Bi-CGSTAB Bi-conjugate gradient stabilized
CP CANDECOMP/PARAFAC
CPU Central processing unit
FEM Finite element method
GMRES Generalized minimal residual
HOPGD Higher-order proper generalized decomposition
HOSVD Higher-order singular value decomposition
IDR Induced dimension reduction
IRKA Iterative rational Krylov algorithm
LTI Linear time-invariant
MIMO Multiple input, multiple output
MINRES Minimal residual
NEP Nonlinear eigenvalue problem
PDE Partial differential equation
PEP Polynomial eigenvalue problem
PGD Proper generalized decomposition
SISO Single input, single output
The hardest to learn was the least complicated

Indigo Girls
Part I: Introduction and Preliminaries
Chapter 1

Introduction

This thesis follows a compilation format and consists of two parts. In Part I, we introduce and summarize the topics covered in the thesis. Chapter 1 provides motivation for the research conducted and offers a summary of related works, aiming to place the appended papers in a broader context. Chapter 2 includes well-established findings relevant to the results presented in the articles, and Chapter 3 summarizes the appended papers and clarifies the author’s contributions. Part II constitutes the main scientific contribution of this thesis and consists of a collection of four papers.

1.1 Linear systems with constant coefficients

We consider approximating the solutions to linear systems described by

$$Ax = b,$$  \hspace{1cm} (1.1)

where $A \in \mathbb{R}^{n \times n}$ is nonsingular, $x \in \mathbb{R}^n$, and $b \in \mathbb{R}^n$. Linear systems of equations with symmetric and nonsymmetric system matrices $A$ arise in the context of engineering and scientific problems. These problems are also present within applications of mathematics to problems in the social sciences, business, and economics [63]. For example, such systems appear when computing the approximate solution to partial differential equations (PDEs) via finite differences or the finite element method (FEM) [48], as well as within finite volume method discretizations for modeling time-dependent PDEs [36].

Linear systems also appear in the analysis and modeling of electrical circuits, where nonlinear systems are approximated by linear ones [11] and in structural analysis, when one is interested in determining the effect of loads on physical structures.
and their components [25]. The possibility to solve larger linear systems allows, for example, a finer FEM-discretization and, therefore, higher accuracy and better modeling possibilities of physical systems [52].

The methods proposed in this thesis can be used to solve linear systems with complex coefficients. The formulation in (1.1) was chosen because many of the theoretical results in the appended papers are based on the assumption that the systems are real-valued. However, the algorithms developed in Paper A and some simulations shown in Paper B involve complex linear systems. In particular, this thesis develops methods for solving linear systems with nonsymmetric system matrices.

**Iterative vs direct methods**

There are two main classes of methods to solve the linear system described by (1.1): direct methods and iterative methods. For linear systems of the form (1.1), direct methods require $O(n^3)$ work. This is too expensive for the applications in this thesis, considering that the input is usually only $O(n^2)$, i.e., typically a vector paired with a sparse matrix. An example of such a method to solve (1.1) is the well-known Gaussian elimination method. This approach requires $O(n)$ steps, where each step requires $O(n^2)$ work. If Gaussian elimination is performed on a computer, it is typically only accurate to machine precision. Figure 1.1 is from [59]. This visualization shows that direct algorithms typically do not decrease the residual until $O(n^3)$ operations have been performed. In other words, the solutions obtained through direct methods are only valuable once the entire, computationally demanding, algorithm has been executed.

On the other hand, the work per step of iterative methods depends on the structure of the matrix. Specifically, each step of the algorithm requires $x \mapsto Ax$, and some iterative algorithms also need the computation $x \mapsto A^T x$. These multiplications are the most costly part of each step, but can be carried out in an efficient way if, for example, $A$ is very sparse. Iterative methods can, in the worst case scenario, require $O(n)$ steps before terminating, where each step requires $O(n^2)$ work. In certain circumstances, iterative methods achieve convergence in $O(1)$ steps, and the work per step can be as low as $O(n)$. This results in a method with total cost of $O(n)$. Typically, however, the total cost of solving linear systems of the form (1.1) with iterative methods is $O(n^2)$. Notably, iterative methods can, like direct methods, achieve accuracy to machine precision. Figure 1.1 shows that, under good conditions, iterative algorithms decrease the residual on each step [59].

The convergence of iterative methods can stagnate, and the prevention of this issue remains an area of active research; see the section on preconditioning, beginning on page 8. Papers A, B, and C of this thesis develop iterative methods for solving large, sparse linear systems, where the system matrix is nonsymmetric. Paper D makes use of the methods proposed in Paper C to obtain solutions to linear systems, then builds a reduced order model from the sampled solutions.
Krylov subspace methods

Krylov subspace methods are among the most important iterative techniques available for solving large linear systems of the form (1.1). Methods of this type are based on a projection onto Krylov subspaces, where the projection can be orthogonal or oblique. The Krylov subspace of dimension $k \in \mathbb{N}$ associated with the matrix $A \in \mathbb{R}^{n \times n}$ and the vector $b \in \mathbb{R}^n$ is defined by

$$K_k(A, b) := \text{span}\{b, Ab, A^2b, \ldots, A^{k-1}b\}. \quad (1.2)$$

In this way, Krylov subspaces are spanned by vectors of the form $p(A)b$, where $p$ is a polynomial [48].

There are many different Krylov subspace methods for solving linear systems, and, depending on context, certain methods are more suitable than others. For example, some methods are preferable when the system matrix is symmetric, while for others, the convergence is characterized by the eigenvalues of $A$ [59]. In general, Krylov subspace methods do not require direct access to the system matrix in order to approximate the solution to (1.1). Instead, one needs only a procedure to compute matrix-vector products with $A$ and for some Krylov subspace methods with $A^T$. This is an attractive feature, particularly in cases where the system matrix is large.

The algorithms proposed in Papers A, B, and C to solve large, sparse, nonsymmetric linear systems are Krylov subspace methods. We describe two foundational Krylov subspace methods below. The novel algorithms proposed in this thesis were inspired by these methods and developed to inherit their robustness, accuracy, and efficiency.

GMRES for linear systems

The generalized minimal residual method (GMRES) [50] is one of the most popular iterative methods for solving linear systems with nonsymmetric system matrices [52]...
and is a generalization of the minimal residual method (MINRES) for symmetric, indefinite matrices [45]. We briefly summarize the approach of GMRES as follows.

Let \( x_k \) denote the \( k \)th iterate produced by GMRES applied to the linear system (1.1), and, without loss of generality, assume \( x_0 = 0 \). The approximation \( x_k \) is the vector in \( \mathcal{K}_k(A, b) \) that minimizes the norm of the \( k \)th residual, i.e.,

\[
x_k := \arg\min_{x \in \mathcal{K}_k(A, b)} \| b - Ax \|_2.
\]

The GMRES algorithm uses an Arnoldi process to generate an orthogonal basis for the Krylov subspace. After \( k \) iterations of the Arnoldi method, the relations

\[
AQ_k = Q_{k+1}H_k
\]

hold, where the orthonormal columns of \( Q_k \in \mathbb{R}^{n \times k} \) span \( \mathcal{K}_k(A, b) \) and \( H_k \in \mathbb{R}^{k+1 \times k} \) is an upper Hessenberg matrix, i.e., \( (H_k)_{i,j} = 0 \) for \( i > j + 1 \). The iterate \( x_k \) is computed as

\[
x_k = Q_kz_k,
\]

where \( z_k = \arg\min_{z \in \mathbb{R}^k} \| H_k z - \beta e_1 \|_2 \) with \( \beta := \| b \| \) and \( e_1 = [1 \ 0 \ldots \ 0]^T \in \mathbb{R}^{k+1} \). Note, the full basis matrix \( Q_k \) must be stored.

In particular, the choice to orthogonalize the basis vectors stems from the observation that the vectors \( A^i b \) tend to become less independent as \( i \in \mathbb{N} \) increases. An orthogonalization process, typically based on a Gram–Schmidt or Householder procedure, alleviates this problem and leads to a more stable method [52]. Furthermore, the norm of the residual vectors produced by the iterates of GMRES are always nonincreasing, i.e.,

\[
\| r_k \| \leq \| r_{k-1} \| \leq \ldots \leq \| r_0 \|,
\]

since \( \mathcal{K}_1(A, b) \subseteq \ldots \subseteq \mathcal{K}_k(A, b) \subseteq \mathcal{K}_{k+1}(A, b) \) and \( r_i \in \mathcal{K}_{i+1}(A, b) \). Moreover, \( \| r_n \| = 0 \), ignoring rounding errors. This is because \( \text{span}\{q_1, \ldots, q_n\} = \mathbb{R}^n \).

The convergence of GMRES is often rapid [59]. The GMRES method applied to the linear system (1.1) is shown in Algorithm 1, and additional theoretical results concerning GMRES can be found in Chapter 2.

### Bi-CG for linear systems

The bi-conjugate gradient method (Bi-CG) [17, 35] is another fundamental Krylov subspace method for solving linear systems with nonsymmetric system matrices. When used to solve systems of the form (1.1), this method requires the construction of a basis for the Krylov subspace \( \mathcal{K}_k(A, b) \), as well as one for the Krylov subspace defined by

\[
\mathcal{K}_k(A^T, c) := \text{span}\{c, A^T c, (A^T)^2 c, \ldots, (A^T)^{k-1} c\}, \tag{1.5}
\]
Algorithm 1: GMRES

**input:** $A \in \mathbb{R}^{n \times n}$, $b \in \mathbb{R}^n$, $k \in \mathbb{N}$

**output:** Approximate solution $x_k$ to $Ax = b$ from Krylov subspace of dimension $k$

1. $\beta := \|b\|$
2. $q_1 = b/\beta \in \mathbb{R}^n$
3. $Q_1 = [q_1] \in \mathbb{R}^{n \times 1}$
4. $H_0 = \text{empty matrix}$
5. for $i = 1, 2, \ldots, k$ do
   6. $y = Aq_i$
   7. Orthogonalize $y$ against $q_1, \ldots, q_i$ by a Gram–Schmidt process:
      - $h_i = Q_i^T y$
      - $y_\perp = y - Q_i h_i$
   8. Possibly repeat Step 7
   9. Compute $\beta_i = \|y_\perp\|$
10. Let $q_{i+1} = y_\perp / \beta_i$
11. Let $Q_{i+1} = [Q_i, q_{i+1}] \in \mathbb{R}^{n \times (i+1)}$
12. Let $H_i = \begin{bmatrix} H_{i-1} & h_i \\ \beta_i & \end{bmatrix} \in \mathbb{R}^{(i+1) \times i}$
13. end
14. Compute $z_k = \text{argmin}_{z \in \mathbb{R}^k} \|H_k z - \beta e_1\|_2$
15. Return $x_k = Q_k z_k$

where $c \in \mathbb{R}^n$ and $c^T b \neq 0$. This is done by a Lanczos biorthogonalization procedure. Constructing the basis for (1.5) requires the action of the matrix $A^T$, which, in some applications, may not be available. We briefly summarize the Bi-CG method as follows.

As a result of the biorthogonalization process, the relation

$$T_k = W_k^T A V_k$$

holds. Specifically, $T_k \in \mathbb{R}^{k \times k}$ is tridiagonal and the columns of $V_k \in \mathbb{R}^{n \times k}$ are a basis for $\mathcal{K}_k(A, b)$, where $v_1 = b/\beta$ and $\beta := \|b\|$. Similarly, the columns of $W_k \in \mathbb{R}^{n \times k}$ span (1.5) and $W_k^T V_k = V_k^T W_k = I_k$. Let $x_k$ denote the $k$th iterate of Bi-CG applied to (1.1). Here, we assume $x_0 := 0$ without loss of generality, and the $k$th iterate is defined as

$$x_k = V_k T_k^{-1} (\beta e_1),$$

where $e_1 = [1 \ 0 \ \ldots \ 0]^T \in \mathbb{R}^k$.

The nonincreasing residual norm property (1.4) does not always hold with Bi-CG. While GMRES chooses the iterate in $\mathcal{K}_k(A, b)$ that minimizes the residual norm, Bi-CG does not, despite operating within this subspace. The convergence of Bi-CG can be slow and erratic in practice [59].
Bi-CG is a three-term recurrence method. In particular, $v_i$ is computed by a linear combination of $v_{i-1}$, $v_{i-2}$, and $Av_{i-1}$. An analogous relation holds for the basis vectors of (1.5). The full basis matrices do not need to be stored, and the Bi-CG method has lower storage and memory requirements than GMRES does. This property has particular significance when solving linear systems that take many iterations to converge.

The standard implementation of the Bi-CG algorithm applied to the linear system (1.1) is shown in Algorithm 2, and this method also returns an approximation to the so-called dual system given by $A^T z = c$. This implementation is based on an implicit LU factorization of the matrix $T_k$; see [48] for more details. Note, the Bi-CG method carries the risk of breakdown due to the underlying Lanczos biorthogonalization process [26]. In particular, this occurs when $s_i^T r_i = 0$ on line 8 of the algorithm; see Chapter 2 for a more rigorous explanation.

**Algorithm 2: Bi-conjugate gradient method**

```plaintext
input: $A \in \mathbb{R}^{n \times n}$, $b \in \mathbb{R}^n$, $c \in \mathbb{R}^n$ such that $c^T b \neq 0$, $k \in \mathbb{N}$
output: Approximate solutions $x_k$ of $Ax = b$ and $z_k$ of $A^T z = c$ from Krylov subspace of dimension $k$

1. $x_0 = 0$, $z_0 = 0$, $p_0 = r_0 = b$, $q_0 = s_0 = c$

2. for $i = 1, 2, \ldots, k$
   3. $\alpha_i = (s_{i-1}^T r_{i-1})/(q_{i-1}^T A p_{i-1})$
   4. $x_i = x_{i-1} + \alpha_i p_{i-1}$
   5. $z_i = z_{i-1} + \alpha_i q_{i-1}$
   6. $r_i = r_{i-1} - \alpha_i A p_{i-1}$
   7. $s_i = s_{i-1} - \overline{\alpha}_i A^T q_{i-1}$
   8. $\beta_i = (s_i^T r_i)/(s_{i-1}^T r_{i-1})$
   9. $p_i = r_i + \beta_i p_{i-1}$
10. $q_i = s_i + \overline{\beta}_i q_{i-1}$

end

12. Return $x_k, z_k$
```

**Preconditioning for Krylov subspace methods**

Utilizing preconditioning in the context of Krylov subspace methods can lead to methods that are more efficient overall. The idea behind preconditioning in Krylov subspace methods is to solve a transformed system which has the same solution as the original problem but is likely to be easier to solve with an iterative method. Preconditioning is only suitable when it can be performed efficiently.

More specifically, the first step is to find a nonsingular preconditioning matrix $M \in \mathbb{R}^{n \times n}$ which is, in some way, close to $A$. The matrix $M$ should also be chosen such that linear systems with this system matrix are inexpensive to solve. This
is an important requirement, since preconditioned Krylov subspace methods must perform an additional solve with the preconditioning matrix at each step [48].

There are three standard ways to apply the preconditioner when solving the linear system (1.1). One choice is to apply the preconditioner from the right, and, as a result, solve the systems

$$AM^{-1}u = b, \quad x = M^{-1}u.$$  

Another option is to consider left preconditioning, i.e., approximating the transformed system

$$M^{-1}Ax = M^{-1}b.$$  

Note, we do not form the matrix $AM^{-1}$ or $M^{-1}A$, but rather rely on methods to perform a linear solve with $M$ at each step of the Krylov algorithm. This can be done, for example, via an iterative method. A final choice is to consider

$$M_L^{-1}AM_R^{-1}u = M_L^{-1}b, \quad x = M_R^{-1}u,$$

which can be used when $M$ is available in the factored form $M = M_LM_R$. In this case, the matrices $M_L$ are $M_R$ are typically triangular [48].

Preconditioning is an area of active research. Papers B, C, and D consider Krylov subspace methods with preconditioning for solving a discretized Helmholtz equation. Specifically, right preconditioning is used, though, in certain circumstances, left preconditioning could have been considered. See also, for instance, the works [13, 14, 15, 49] for other successful approaches utilizing preconditioning in the context of solving a discretized Helmholtz equation.

We summarize below two Krylov subspace methods that apply different preconditioners at each iteration. The methods developed in Papers B and C are inspired by this framework.

**Flexible GMRES**

The prior work [47] is based on the GMRES method with right preconditioning. Instead of applying the preconditioner $M$ at each step, the preconditioner in this algorithm is permitted to change from iteration to iteration. Concretely, a linear system with system matrix $M_i$ must be solved at the $i$th iteration, and the method is referred to as flexible GMRES.

Unlike iterates from the standard GMRES method, approximations obtained from flexible GMRES do not come from a Krylov subspace; see [54] for details. Additionally, the flexible GMRES procedure requires the storage of two basis matrices, of which one is the basis for a Krylov subspace. In contrast, the original GMRES algorithm, applied with or without preconditioning, requires only the storage of a single basis matrix for a Krylov subspace. Experimentally, using flexible GMRES can be worth the extra cost. See, for example, [9, 10, 24], for works based on its success.
Inexact Krylov subspace methods

In [5], [6], and [7], the authors experimented with the accuracy of the matrix-vector multiplications \( x \mapsto Ax \) in Krylov subspace methods. When approximating a system of the form (1.1), it was observed that the norm of the error in the multiplication could increase as the norm of the outer residual, i.e., \( r_i = b - Ax_i \), decreased, and this could be done without destroying the convergence of the iterative method. Later, in [54], the authors provided rigorous theory describing how inexactly the matrix-vector products could be carried out while maintaining convergence for many fundamental Krylov subspace methods. Moreover, the authors provided computable bounds on the residual of these inexact Krylov subspace methods.

Papers B and C use the theory of inexact Krylov methods in order to apply a preconditioner in a more efficient way. Specifically, the preconditioners are applied using iterative solvers, and the tolerance is relaxed as the outer Krylov iteration proceeds. In this way, the approach is also similar to the framework proposed in flexible GMRES [47], and two basis matrices need to be stored. Similar approaches have been taken in, for example, the prior works [34] and [62].

1.2 Parameterized linear systems

Specifically, the research performed in this thesis considers parameterized linear systems of the form

\[
A(\mu)x(\mu) = b, \quad \mu = [\mu_1, \ldots, \mu_\ell]^T \in \mathbb{R}^\ell,
\]

for nonsingular \( A(\mu) \in \mathbb{R}^{n \times n} \), \( x(\mu) \in \mathbb{R}^n \), and \( b \in \mathbb{R}^n \). The novel methods we propose are developed to approximate solutions \( x(\mu) \) for many different values of \( \mu \). The subsequent sections provide descriptions of scientific applications that require solutions to linear systems of this form. The purpose of this is to motivate the work that was carried out and to clarify the importance of our results. An overview of how such systems are solved in the appended papers is also provided.

Parameterized shifted linear systems

Solving linear systems of the form

\[
(A + \mu I)x(\mu) = b
\]

in a computationally efficient way for many different values of the parameter \( \mu \in \mathbb{R} \) is of scientific interest. Here the system matrix differs from \( A \in \mathbb{R}^{n \times n} \) by a scalar multiple of the identity matrix, and this linear system is of the same form as (1.6) with \( \ell = 1 \). Specifically, the dependence on the parameter is linear, and we refer to such linear systems as parameterized shifted linear systems.
1.2. Parameterized linear systems

Linear systems of the form (1.7) arise, for instance, in the context of model order reduction [27], Tikhonov regularization for ill-posed problems [23, 32], and oscillatory hydraulic tomography [51], used to image the subsurface of the earth. Many prior works have been successfully developed for solving these linear systems; see, for example, [4, 18, 22, 56]. Parameterized shifted linear systems appear throughout this thesis.

Krylov subspace methods are attractive for solving linear systems of this variety. This is because the basis matrix for one Krylov subspace can be used to compute solutions corresponding to many different \( \mu \) simultaneously, due to the shift-invariance property of Krylov subspaces:

\[
\mathcal{K}_k(A, b) = \mathcal{K}_k(A + \mu I, b).
\]

We show how one basis matrix can be used to solve (1.7) for many different \( \mu \) using the GMRES method as follows.

The Arnoldi relations (1.3) can be expressed equivalently as

\[
(A + \mu I)Q_k = Q_{k+1}(H_k + \mu L_k),
\]

where \( L_k \in \mathbb{R}^{k+1 \times k} \) is the identity matrix with an extra row of zeros and \( Q_{k+1}L_k = Q_k \). Furthermore, the columns of \( Q_k \) are an orthonormal basis for \( \mathcal{K}_k(A + \mu I, b) \), and this basis matrix is generated independently of \( \mu \), as described in Algorithm 1. The \( k \)th iterate of the GMRES method applied to (1.7) can be computed as

\[
x_k(\mu) = Q_kz_k(\mu),
\]

where

\[
z_k(\mu) = \arg\min_{z \in \mathbb{R}^k} \| (H_k + \mu L_k)z - \beta e_1 \|_2
\]

and \( e_1 = [1 \ 0 \ \ldots \ 0]^T \in \mathbb{R}^{k+1} \). Thus, after the Arnoldi portion of the method has been executed, solving (1.7) requires only the solution to a small least squares problem and a matrix-vector multiplication for each \( \mu \) of interest. Moreover, \( x_k(\mu) \in \mathcal{K}_k(A + \mu I, b) \).

Papers A, B, and C in this thesis develop computationally effective approaches for solving shifted systems based on a similar framework. Though a GMRES procedure was described here and used in Paper A, other Krylov subspace methods can be used in an analogous way. Specifically, Paper B considers an approach based on inexact GMRES, and the algorithm proposed in Paper C is based on the Bi-CG method.

Parameterized partial differential equations

Consider a discretization of the inhomogeneous, indefinite Helmholtz equation given by

\[
(\nabla^2 + \kappa(\mu, x)^2) u(x) = h(x), \quad x \in \Omega
\]

(1.9)
with a homogeneous Dirichlet boundary condition: \( u(x) = 0 \), for \( x \in \partial \Omega \). Here, \( \kappa \) is referred to as the wave number, and we assume the function \( \kappa(\mu, x)^2 > 0 \) is separable in \( \mu \) and \( x \), i.e.,

\[
\kappa(\mu, x)^2 := f_1(\mu)g_1(x) + \ldots + f_{nf}(\mu)g_{nf}(x),
\]

where \( \mu \in \mathbb{R}^\ell, f_i : \mathbb{R}^\ell \to \mathbb{R}, g_i : \mathbb{R}^d \to \mathbb{R}, \) and \( d \) is the dimension of the spatial domain \( \Omega \). Furthermore, \( h \) represents sources inside the domain, and the solution \( u(x) \) describes wave propagation [16]. This elliptic PDE was chosen as a representative example. We refer to \( \mu \) as the input parameter. Observe that a change in the parameter \( \mu \) leads to a change in the response, i.e., the solution \( u(x) \). In this way, a discretization of this problem is of the form (1.6).

Parameterized PDEs similar to (1.9) arise in the study of geophysics; see, for instance, [65]. Specifically, in the context of the parameterized Helmholtz equation, the input parameter can represent the angular frequency of the wave or the direction of the incoming wave [41]. Alternatively, \( \mu \) can describe the geometry of the domain or the materials through which the waves move [53].

The algorithms proposed in this thesis can be seen as tools for uncertainty quantification. More concretely, when there is uncertainty in the input parameter, \( \mu \) can be treated as a random variable. For each observation of \( \mu \), we can obtain a corresponding realization of the solution \( u(x) \), which is also a random variable. This is referred to as forward propagation of uncertainty. Our methods offer a computationally efficient way to approximate \( u(x) \) for many values of the parameter. Having access to many such approximations allows us to compute descriptive properties, such as the mean and variance, of the solutions \( u(x) \) [57].

**Numerical examples**

The numerical simulations in Papers A, B, C, and D consider a discretization of (1.9) by the finite element method, though other discretization methods could have been used. The resulting discretized system is of the form (1.6), where

\[
A(\mu) := A_0 + f_1(\mu)A_1 + \ldots + f_{nf}(\mu)A_{nf}
\]

and \( b \) is the corresponding load vector. Specifically, \( A_0, \ldots, A_{nf} \) are large and sparse, the functions \( f_i \) are nonlinear in \( \mu \), and we generate the discretized system using FEniCS [2]. The solution to the discretized problem is an approximation to the Helmholtz equation (1.9). Performing the discretization with a finer mesh leads to a more accurate model of the PDE and a larger linear system [52]. Thus, developing efficient methods for large, sparse systems is of interest.

Furthermore, we are interested in the solutions for many different \( \mu \). The obvious approach of directly interpolating a few approximations is not a suitable strategy because solutions to the parameterized Helmholtz equation are extremely sensitive to the parameters [41]. As an alternative, Papers A, B, and C develop
1.2. Parameterized linear systems

Krylov subspace methods to obtain approximations corresponding to many different values of a scalar $\mu$. Paper D proposes a reduced order model to achieve the same goal with $\mu \in \mathbb{R}^2$.

More precisely, Papers A, B, and C are based on solving an equivalent linearized system of the form (1.7) with Krylov subspace methods. Though using a linearization requires the solution to a larger system, sparsity patterns keep the cost low, particularly with respect to matrix-vector multiplications. Moreover, approximations to (1.9) corresponding to many different scalar $\mu$ can be generated from one basis matrix; see the section on solving parameterized shifted linear systems with Krylov subspace methods, beginning on page 10. The algorithm developed in Paper D relies on sampling performed using a modified version of the method proposed in Paper C.

**Outlook**

The numerical solution to (1.9) is usually difficult to obtain with standard iterative methods, especially if the wave number $\kappa$ is large. Improvements can be made by, for example, considering Krylov subspace methods with preconditioning by incomplete factorizations, domain decomposition, or multigrid methods. These approaches lead to better converge of the iterative method, though the number of iterations tends to increase with the size of the wave number. Building algorithms whose convergence is independent of the wave number remains a difficult problem, and sophisticated solvers are required [16]. In this thesis, we consider numerical simulations of (1.9) for modestly sized wave numbers without tackling the problem of increasing wave numbers. Continued research into combining our approaches with preconditioning designed for the Helmholtz equation with large wave numbers could lead to promising approaches.

**Linear systems with time delays**

Parameterized linear systems of the form (1.6) also arise in the context of linear time-invariant (LTI) systems. Specifically, we consider systems with time-delays. These systems are applicable in various domains, including flight control, chemical/process control, and communication networks. The delays can be in the state, the control input, or the measurements [20, 40].

In particular, we develop methods to study the system described by

\[ \dot{x}(t) = A_0 x(t) + A_1 x(t - \tau) - b v(t), \]  
\[ y(t) = C^T x(t). \]  

Here, the matrices $A_0, A_1 \in \mathbb{R}^{n \times n}$ have random entries, $b \in \mathbb{R}^n$ is a random vector, and $C = I \in \mathbb{R}^{n \times n}$. The delay is given by $\tau \in \mathbb{R}$, where $\tau > 0$, $v(t) \in \mathbb{R}$ describes the input, and $y(t) \in \mathbb{R}^n$ is the output. Additionally, $x(t) \in \mathbb{R}^n$ represents the state vector at time $t \in \mathbb{R}$, and $\dot{x}(t) \in \mathbb{R}^n$ describes how the state changes with respect to time. Without loss of generality, we take $\tau = 1$. 

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Numerical examples

Simulations involving the transfer function of the system (1.10) appear in Papers A, B, and C. The transfer function of an LTI system is the ratio of the Laplace transform of the output variable to the Laplace transform of the input variable with all initial conditions set to zero. Specifically, the transfer function is parameterized in the Laplace variable and describes the input-output behavior of the system without giving any information about the system’s internal structure. The transfer function is independent of the input, and, from analyzing the poles of the transfer function, the stability of the system can be determined [43]. In particular, the transfer function of the system (1.10) is the solution to the parameterized linear system (1.6), where

\[ A(\mu) := -\mu I + A_0 + A_1 e^{-\mu}; \]  

see [66] for details. Here, the parameter \( \mu \) is the Laplace variable, and we are interested in evaluating the transfer function for many different \( \mu \).

The algorithms developed in this thesis consider Krylov subspace methods to solve an equivalent linearized problem of the form (1.7) in an efficient way. This procedure involves solving a linear system which is larger than the original one. Sparsity patterns in the linearization are exploited in order to develop computationally efficient methods. Further, solutions corresponding to many different values of the parameter \( \mu \) are generated from just one basis matrix; see the section on parameterized shifted linear, beginning on page 10.

Outlook

Improved modeling of transfer functions leads to better understanding of engineering control systems in general [44]. Particularly, in robust control of linear systems, stability and performance are often evaluated using \( \mathcal{H}_\infty \)-norms of transfer functions [64]. For systems with transfer function \( H \), this norm is defined as

\[ \| H \|_{\mathcal{H}_\infty} := \sup_{\omega \in \mathbb{R}} |H(i\omega)|, \]  

for stable single input, single output systems (SISO),

\[ \| H \|_{\mathcal{H}_\infty} := \sup_{\omega \in \mathbb{R}} \sigma_{\max}(H(i\omega)), \]  

(1.12)

for stable multiple input, multiple output systems (MIMO), and \( \| H \|_{\mathcal{H}_\infty} := \infty \) for unstable systems. The quantity in (1.12) can be interpreted as the largest singular value of the transfer function \( H \) over the imaginary axis. Many methods have been proposed to approximate the \( \mathcal{H}_\infty \)-norm without evaluating the transfer function directly; see, for instance, the approaches developed in [8, 28, 29]. Clearly, if the dimension \( n \) is large, evaluating the transfer function for even a few different values of \( \omega \) is computationally demanding.
Furthermore, the $H_2$-norm is important as a robustness measure with respect to noise or external disturbances [46]. We define the $H_2$-norm of a transfer function $H$ by

$$\|H\|_{H_2} := \left(\frac{1}{2\pi} \int_{-\infty}^{\infty} |H(i\omega)|^2 \, d\omega \right)^{\frac{1}{2}},$$

for stable SISO systems, and

$$\|H\|_{H_2} := \left(\frac{1}{2\pi} \int_{-\infty}^{\infty} \text{tr}(H(-i\omega)^T H(i\omega)) \, d\omega \right)^{\frac{1}{2}},$$

for stable MIMO systems. There are many successful methods for LTI systems which utilize the $H_2$-norm. Notably, in [27], the iterative rational Krylov algorithm (IRKA) was proposed for producing the reduced order model with transfer function $H_r$ that minimizes the quantity $\|H - H_r\|_{H_2}$ in the context of SISO LTI systems without delays. Specifically, the method obtains the best stable dynamical system of dimension $r$ for approximating the original system, where $r \ll n$.

Efficient and accurate methods for evaluating transfer functions for many different values of the parameter could lead to promising new methods for computing $H_\infty$-norms and $H_2$-norms. Such results could lead to improved methods for evaluating and modeling control systems.

### Parameter estimation problems for PDEs

The goal of parameter estimation in the context of the parameterized linear system (1.6) is to determine the value of the unknown parameter $\mu$ that best fits the corresponding observed solution $x(\mu)$. The parameter estimation problem is an inverse problem, as we aim to infer properties of the system that produced a given solution [58]. Inverse problems are typically ill-posed. Specifically, there may be no solution, the solution may not be unique, or there may be a unique solution that is very sensitive to changes in the response [57].

Consider the problem of having a set of solutions to the system (1.6), and let these solutions be denoted $\{x_1, \ldots, x_N\}$, for $x_i \in \mathbb{R}^n$. Here the corresponding parameters $\mu_i$ are unknown and of interest, and the observed solutions are noisy, i.e., $x_i \simeq x(\mu_i) + \varepsilon_i$, where $\varepsilon_i \sim \mathcal{N}(0, \sigma^2)$ is assumed. In this context, $\mu_i$ is an observation of the random variable $\mu$, and, thus, $x_i$ is the corresponding realization of the random variable $x$. The noise in the observations could arise, for example, from measurement error.

### Numerical examples

The algorithms developed in this thesis return functions $\tilde{x}(\mu) \in \mathbb{R}^n$ such that $\tilde{x}(\mu) \approx x(\mu)$, and the functions can be evaluated in a computationally efficient way for many values of the parameter. As a result, the solution of a parameter estimation problem
can be obtained by a direct application of the proposed algorithms. Concretely, we obtain an approximation to $\mu$ by considering the minimizations

$$
\mu_i^* = \underset{\mu}{\text{argmin}} \|x_i - \tilde{x}(\mu)\|_2, \quad i = 1, \ldots, N.
$$

(1.13)

Numerical simulations involving a parameter estimation in the context of PDEs can be found in Paper D with $\mu = [\mu_1, \mu_2]^T \in \mathbb{R}^2$. In one simulation, $\mu$ describes unknown material parameters in the context of a Helmholtz equation. A second simulation describes an advection-diffusion equation, which can be used to model particle transport; see [12, 60]. Here, the unknown parameters correspond to an advection parameter as well as a diffusion coefficient. It is assumed that the solutions in these examples depend on $\mu$ in a nonlinear way.

One could, alternatively, approach this parameter estimation problem by solving (1.6) for many different values of the parameter $\mu$ using standard methods and then compare each $x_i$ to the collection of computed solutions. This becomes infeasible when many precomputed solutions to large systems are required or the dimension of $\mu \in \mathbb{R}^\ell$ grows.

**Outlook**

From many observations of solutions to a parameterized linear system, we can determine a posterior distribution for the parameter $\mu$. In other words, given $x_i$, we compute $\mu_i^*$ and construct a probability for $\mu$ using this information combined with a prior distribution of $\mu$. Specifically, this probability is based on the realizations $x_i$ [57].

Parameter estimation problems can, for example, help oil companies locate oil deposits to make a profit or help seismologists make predictions regarding earthquakes [57]. Other parameter estimation problems in the context of parameterized PDEs have been tackled in, for instance, [38, 42, 53]. See also [39] and references therein for a thorough summary of both theoretical and numerical methods for parameter estimation in the context of parameterized Helmholtz equations. Improved methods for parameterized linear systems could lead to better methods for tackling these types of problems.

**Parameterized right-hand sides**

The linear system (1.6) is the focus of the scientific work presented in this thesis. Thus, we consider only parameterized linear systems with constant right-hand sides. The prior work [55] studied the set of parameterized linear systems of dimension $n \times n$ with complex coefficients given by

$$
(A + \mu_i I)x(\mu_i) = b(\mu_i), \quad \mu_i \in \mathbb{C}, \quad \text{for} \quad i = 1, \ldots, N.
$$

(1.14)

In that setting, the right-hand side vectors $b(\mu_i)$ are unrelated, and the systems were solved simultaneously by block Krylov subspace methods. The general approach
could be adapted to our setting to, for example, solve parameterized PDEs with
certain types of parameter-dependent boundary conditions. However, this would
require additional research. The method presented in [55] is briefly summarized as
follows.

Without loss of generality, let \( x_0 := 0 \in \mathbb{R}^n \) be the initial approximation to each
system in (1.14), and let the initial residuals be given by \( r_0(\mu_i) := b(\mu_i) - A x_0 \). The
shift-invariance of Krylov subspaces described by (1.8) cannot be exploited in this
setting. This is because the vectors \( b(\mu_i) \) are linearly independent and, thus,

\[
\mathcal{K}_k(A + \mu_i I, b(\mu_i)) \neq \mathcal{K}_k(A + \mu_j I, b(\mu_j)), \quad \text{for} \quad i \neq j.
\]

We focus instead on block Krylov subspace methods. Consider the Sylvester
equation

\[
AX^\mu + X^\mu D = B^\mu, \quad (1.15)
\]

where \( D := \text{diag}(\mu_1, \ldots, \mu_N) \in \mathbb{C}^{N \times N} \), \( B^\mu := [b(\mu_1) \ldots b(\mu_N)] \in \mathbb{C}^{n \times N} \) are given
and the matrix

\[
X^\mu := [x(\mu_1) \ldots x(\mu_N)] \in \mathbb{C}^{n \times N}
\]

is unknown. In this way, solutions to (1.14) and (1.15) are equivalent. Let \( X_0^\mu :=
0 \in \mathbb{C}^{n \times N} \) be the initial approximation to \( X^\mu \) and

\[
R_0^\mu := B^\mu - AX_0^\mu - X_0^\mu D = B^\mu
\]

be the corresponding initial residual. The block Krylov subspace generated from \( A \)
and \( B^\mu \) is defined as

\[
\mathbb{K}_k(A, B^\mu) := \mathcal{K}_k(A, b(\mu_1)) + \mathcal{K}_k(A, b(\mu_2)) + \ldots + \mathcal{K}_k(A, b(\mu_N)),
\]

and the relation

\[
\mathbb{K}_k(A, B^\mu) = \mathbb{K}_k(A + \mu I, B^\mu), \quad \mu \in \mathbb{C},
\]

is referred to as the shift-invariance property of block Krylov subspaces. It is assumed
that \( A \) and \( B^\mu \) are chosen such that \( \dim(\mathbb{K}_k(A, B^\mu)) = kN \).

The subspace \( \mathbb{K}_k(A, B^\mu) \) can be represented in terms of the orthonormal block
Arnoldi basis given by \( \{U_1, \ldots, U_k\} \). Here each \( U_i \in \mathbb{C}^{n \times N} \) has orthonormal columns,
and the columns of \( U_i \) are orthogonal to the columns of \( U_j \), for \( i \neq j \). We define \( U_k :=
\left[ U_1 \ldots U_k \right] \in \mathbb{C}^{n \times kN} \) and denote the block upper Hessenberg matrix generated by
the block Arnoldi method as \( \mathbf{H}_k \in \mathbb{C}^{(k+1)N \times kN} \). The relation

\[
A U_k = U_{k+1} \mathbf{H}_k
\]

holds, as well as the relations

\[
(A + \mu_i I) U_k = U_{k+1} \mathbf{H}_k(\mu_i), \quad (1.16)
\]
where

\[ H_k(\mu_i) = H_k + \mu_i \begin{bmatrix} I_{kN \times kN} & \end{bmatrix}. \]

The shifted Arnoldi relation (1.16) is used to derive Algorithm 3.1 in [55]. Specifically, approximations to the parameterized linear systems (1.14) are elements of the block Krylov subspace \( \mathbb{K}_k(A, B^\mu) \), and the corresponding basis matrix \( U_k \) is used to generate the solutions. Computing each \( x(\mu_i) \) requires the solution to a linear least squares problem and a matrix-vector product. This procedure is similar to the GMRES method applied to the parameterized shifted linear system (1.7), described in the section on shifted parameterized linear systems beginning on page 10.
Chapter 2

Preliminaries

This chapter contains an overview of some well-known concepts and theoretical results from the literature. Specifically, it includes fundamental results related to the Krylov subspace methods that inspired the appended papers. The purpose of this is to provide an overview and foundation for the topics included in this thesis. Since all the results in this chapter are well-established, proofs are omitted. Further details can be found in the books in the reference list at the end of this part.

In particular, the following results pertain to the linear system with constant coefficients described by (1.1) or, equivalently, the parameterized linear system (1.6) evaluated at a specific value of the parameter. We begin with some results regarding matrix factorizations, which are used both in computation and theoretical analysis throughout this dissertation.

**Theorem 2.1** (LU factorization). Let $A \in \mathbb{R}^{n \times n}$ be such that

$$\det(A(1 : i, 1 : i)) \neq 0, \quad i = 1, \ldots, n - 1.$$

Then there exists an LU factorization of $A$, i.e., a lower triangular matrix $L \in \mathbb{R}^{n \times n}$ and an upper triangular matrix $U \in \mathbb{R}^{n \times n}$ such that $A = LU$.

The LU factorization plays an important role in various computational tasks, such as solving systems of linear equations, computing determinants, and inverting matrices. In particular, systems involving lower or upper triangular matrices can be solved efficiently by considering forward or back substitution, respectively.

**Theorem 2.2** (QR factorization). If $A \in \mathbb{R}^{n \times m}$, then there exists an orthogonal matrix $Q \in \mathbb{R}^{n \times n}$ and an upper triangular matrix $R \in \mathbb{R}^{n \times m}$ such that

$$A = QR.$$
Chapter 2. Preliminaries

Since the transpose of an orthogonal matrix is equal to its inverse, the QR factorization helps to simplify many computations. The following theorem is applicable in the case of overdetermined systems.

**Theorem 2.3** (Reduced QR factorization). If $A \in \mathbb{R}^{n \times m}$ has full column rank and $n \geq m$, then the reduced QR factorization of $A$ is given by $Q_1R_1$, where

$$A = QR = \begin{bmatrix} Q_1 & Q_2 \end{bmatrix} \begin{bmatrix} R_1 & 0 \\ 0 & 0 \end{bmatrix} = Q_1R_1. \tag{2.1}$$

Here $R_1 \in \mathbb{R}^{m \times m}$ is an upper triangular matrix, $0$ is a zero matrix of dimension $(n - m) \times m$, $Q_1 \in \mathbb{R}^{n \times m}$, $Q_2 \in \mathbb{R}^{n \times (n - m)}$, and $Q_1$ and $Q_2$ have orthogonal columns. Moreover, the factorization $A = Q_1R_1$ where $R_1$ has positive diagonal entries is unique.

As the algorithms developed in this thesis are Krylov subspace methods, the following definitions are foundational.

**Definition 2.4** (Krylov subspace). The Krylov subspace of dimension $k$ generated by the matrix $A \in \mathbb{R}^{n \times n}$ and the vector $b \in \mathbb{R}^n$ is given by

$$\mathcal{K}_k(A, b) = \text{span}\{b, Ab, \ldots, A^{k-1}b\}. \tag{2.2}$$

**Definition 2.5** (Krylov matrix). The Krylov matrix $K_k \in \mathbb{R}^{n \times k}$ associated with the Krylov subspace (2.2) is defined by

$$K_k := \begin{bmatrix} b & Ab & A^2b & \ldots & A^{k-1}b \end{bmatrix}. \tag{2.3}$$

The parameterized shifted linear system (1.7) arises throughout this thesis. As a result, the algorithms we develop rely heavily on the following theorem.

**Theorem 2.6** (Shift- and scaling-invariance). Krylov subspaces are shift- and scaling-invariant, i.e.,

$$\mathcal{K}_k(A, b) = \mathcal{K}_k(\mu A, b) = \mathcal{K}_k(A - \mu I, b),$$

for nonzero $\mu \in \mathbb{R}$.

### 2.1 Regarding GMRES

The particular Krylov subspace methods developed in Papers A and B were inspired by the GMRES method and inherit many of its properties. When the standard GMRES method is applied to (1.1), an Arnoldi procedure constructs a basis matrix for (2.2) with orthogonal columns. The following theorem formalizes this.
2.1. Regarding GMRES

Theorem 2.7 (Arnoldi relation). Suppose the Arnoldi algorithm applied to the matrix $A \in \mathbb{R}^{n \times n}$ with starting vector $q_1 = b/\beta \in \mathbb{R}^n$ and $\beta := \|b\|$ does not break down after $k$ iterations. Then, the Arnoldi relation

$$AQ_k = Q_{k+1}H_k$$

(2.4)

holds, where $Q_{k+1} = [Q_k \ q_{k+1}] \in \mathbb{R}^{n \times k+1}$, $H_k \in \mathbb{R}^{k+1 \times k}$ is an upper Hessenberg matrix, and $H_{k-1} = H_k(1: k, 1: k-1)$. Furthermore, the columns of $Q_k$ are an orthonormal basis for the Krylov subspace $K_k(A, b)$.

We describe the relation between the basis matrix $Q_k$ in (2.4) and the Krylov matrix associated with the subspace (2.2) in the following theorem.

Theorem 2.8. The matrix $Q_k$ in (2.4) is the reduced QR factor of the Krylov matrix (2.3), i.e.,

$$K_k = Q_kR_k.$$ 

When the GMRES method is used to solve (1.1), the Krylov basis vectors are orthogonalized using a Gram–Schmidt procedure, and the full basis matrix must be stored. This can be costly when many iterations are needed to achieve convergence. We illustrate this as follows.

Corollary 2.9. Let $q_1, \ldots, q_k$ denote the columns of $Q_k$. The Krylov vector $q_{k+1}$ in (2.4) fulfills a $(k+1)$-term recurrence between itself and the previous Krylov vectors since

$$Aq_k = h_{1,k}q_1 + h_{2,k}q_2 + \ldots + h_{k+1,k}q_{k+1},$$

and $AQ_k = h_{1,k}q_1 + h_{2,k}q_2 + \ldots + h_{k+1,k}q_{k+1}$.

Proposition 2.10. The cost of orthogonalizing the columns of the Krylov matrix (2.3) with the Gram–Schmidt procedure is $O(nk^2)$.

Definition 2.11 (Iterates of GMRES). Let $x_0 := 0$. The $k$th iterate of the GMRES method applied to (1.1) is given by

$$x_k := Q_kz_k,$$

where $z_k := \arg\min_{z \in \mathbb{R}^k} \|H_kz - \beta e_1\|_2$, $\beta := \|b\|$, $e_1 = [1 \ 0 \ \ldots \ 0]^T \in \mathbb{R}^{k+1}$, and $Q_k$, $H_k$ are as in (2.4).

Remark 2.12 (Solving the least squares problem). The vector $z_k$ in Definition 2.11 can be obtained in an efficient way by considering a QR decomposition of the Hessenberg matrix $H_k$. This factorization can be reused when computing $x_{k+1}$. 

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The residual vector is used to approximate the error when the true solution is not available. The following fundamentals characterize the residual.

**Definition 2.13** (Residual of iterates, GMRES). The residual of the $k$th iterate of the GMRES method applied to the linear system (1.1) is computed efficiently as
\[ \|r_k\|_2 = \|\beta e_1 - H_k z_k\|_2. \]
Furthermore, $r_k \in \mathcal{K}_{k+1}(A, b)$.

**Definition 2.14** (Petrov-Galerkin condition, GMRES). The property $b - Ax_k \perp A\mathcal{K}_k(A, b)$ is referred to as the Petrov-Galerkin condition for the GMRES method.

**Proposition 2.15** (Minimal residual, GMRES). The $k$th iterate of the GMRES method applied to (1.1) is the vector in $\mathcal{K}_k(A, b)$ which minimizes the $k$th residual norm
\[ \|r_k\|_2 = \|b - Ax_k\|_2. \]

The main convergence theorem for the GMRES method is based on describing the iterates using matrix polynomials. We begin with some basic definitions.

**Definition 2.16** (Space of polynomials). The space of polynomials $P_k$ is given by
\[ P_k := \{ \text{polynomials of degree } \leq k \text{ with } p(0) = 1 \}. \] (2.5)

**Definition 2.17** (Polynomial definition of GMRES iterates). The iterates of the GMRES method applied to (1.1) can be expressed as
\[ x_k = q_k(A)b \]
at iteration $k$, where $q_k$ is a polynomial of degree $k - 1$.

The corresponding residual vectors can also be represented using this polynomial interpretation.

**Definition 2.18** (Polynomial definition of residual, GMRES). The residual of the $k$th iterate of GMRES can be described by
\[ r_k = b - Aq_k(A)b = (I - Aq_k(A))b = p_k(A)b \] (2.6)
with $p_k(z) = 1 - zq_k(z)$ and $p_k \in P_k$.

**Proposition 2.19.** At iteration $k$, the GMRES process applied to (1.1) determines the coefficients of $p_k$ such that $\|r_k\|_2$ is minimized.
We can describe the convergence of the GMRES method applied to (1.1) using properties of the system matrix. The following is the main convergence theorem of the GMRES method.

**Theorem 2.20.** Let $A \in \mathbb{R}^{n \times n}$ be nonsingular and diagonalizable with diagonalization $A = V \Lambda V^{-1}$. Then,

$$\frac{\|r_k\|}{\|b\|} \leq \inf_{p_k \in P_k} \|p_k(A)\| \leq \kappa(V) \inf_{p_k \in P_k} \|p_k\|_{\lambda(A)},$$

where $\kappa(\cdot)$ denotes the condition number, $\lambda(A)$ is the spectrum of $A$, and $\|p_k\|_{\lambda(A)} := \sup_{\lambda_i \in \lambda(A)} |p_k(\lambda_i)|$.

The bound described above includes the condition number of $V$. A large condition number can indicate potential linear dependence among the columns.

The following bound, based on a localization of the eigenvalues of $A$, is a direct consequence of Theorem 2.20.

**Corollary 2.21.** Let $A$ fulfill the same conditions as in Theorem 2.20. Furthermore, let all the eigenvalues $\lambda_i \in \lambda(A)$ be enclosed in a disc of radius $r \in \mathbb{R}$ centered at $c \in \mathbb{C}$. Then,

$$\frac{\|r_k\|}{\|b\|} \leq \kappa(V) \left( \frac{r}{|c|} \right)^k. \quad (2.7)$$

Note, the results in Theorem 2.20 and Corollary 2.21 give a sufficient, but not necessary, condition for convergence. Moreover, the bound in (2.7) is not useful if the eigenvalues are centered too near the origin.

Breakdown can occur in a GMRES context, but is not something to be feared. The following remark formalizes this.

**Remark 2.22 (Breakdown, GMRES).** A so-called lucky breakdown occurs if $h_{k+1,k} = 0$ in Corollary 2.9. In this case, the algorithm terminates with $\|r_k\| = 0$, i.e., $x_k$ is an exact solution.

### 2.2 Regarding Bi-CG

The algorithm developed in Paper C and utilized in Paper D is based on the three-term recurrence Krylov subspace method Bi-CG. Consider approximating the linear system (1.1) with this method. Unlike algorithms based on an Arnoldi process, this approach does not require the storage of a basis matrix or a separate orthogonalization process. Fundamental properties of the Lanczos biorthogonalization procedure, upon which Bi-CG is based, are described as follows.

**Theorem 2.23 (Lanczos biorthogonalization).** The Lanczos biorthogonalization generates matrices $V_k, W_k \in \mathbb{R}^{n \times k}, T_k \in \mathbb{R}^{k \times k}$, and $T_k, \bar{T}_k \in \mathbb{R}^{(k+1) \times k}$ such that the
relations

\[ AV_k = V_k T_k + \beta_k v_{k+1} e_k^T = V_{k+1} T_{k}, \quad (2.8a) \]
\[ A^T W_k = W_k T_k^T + \gamma_k w_{k+1} e_k^T = W_{k+1} T_k^T, \quad (2.8b) \]

hold, where \( e_k \) is the \( k \)th column of the identity matrix of dimension \( k \times k \) and \( \beta_k, \gamma_k \in \mathbb{R} \) are entries of \( T_k \) and \( \bar{T}_k \), respectively. Moreover, the columns of \( V_k \) span the subspace \( K_k(A, b) \), and the columns of \( W_k \) span the subspace \( K_k(A^T, c) \), for \( c \in \mathbb{R}^n \), \( c^T b \neq 0 \).

In particular, this biorthogonalization procedure requires the action of the the matrix \( A \) as well as the action of the adjoint \( A^T \), which may not be available in some applications. The following results clarify the output of this process.

**Corollary 2.24** (Tridiagonal matrices, Lanczos biorthogonalization). In the relations (2.8), the square matrix \( T_k \) has the form

\[ T_k := \begin{bmatrix} \alpha_1 & \gamma_1 & & \\ \beta_1 & \ddots & \ddots \\ & \ddots & \ddots & \ddots \\ & & \beta_{k-1} & \alpha_k \end{bmatrix}, \quad (2.9) \]

and the matrices \( T_k \) and \( \bar{T}_k^T \) are given by

\[ T_k := \begin{bmatrix} \alpha_1 & \gamma_1 & & \\ \beta_1 & \ddots & \ddots \\ & \ddots & \ddots & \ddots \\ & & \beta_{k-1} & \alpha_k \end{bmatrix}, \quad \bar{T}_k^T := \begin{bmatrix} \alpha_1 & \beta_1 & & \\ \gamma_1 & \ddots & \ddots \\ & \ddots & \ddots & \ddots \\ & & \gamma_{k-1} & \alpha_k \end{bmatrix}. \quad (2.10) \]

In particular, \( T_k \) is tridiagonal and \( T_k, \bar{T}_k^T \) are tridiagonal, upper Hessenberg matrices.

**Corollary 2.25.** The Lanczos biorthogonalization procedure gives the relation

\[ W_k^T V_k = V_k^T W_k = I_k, \quad (2.11) \]

where \( I_k \in \mathbb{R}^{k \times k} \) is the identity matrix.

The iterates of the Bi-CG method applied to (1.1) are defined as follows.

**Definition 2.26** (Iterates of Bi-CG). Let \( x_0 = 0 \). The \( k \)th iterate of the Bi-CG method applied to (1.1) is defined as

\[ x_k = V_k T_k^{-1}(\beta e_1), \quad (2.12) \]

where the matrices \( V_k \) and \( T_k \) are as in (2.8) with \( v_1 = b/\beta \) and \( \beta := \|b\| \), and \( e_1 = [1 \ 0 \ \ldots \ 0]^T \in \mathbb{R}^k \).
2.2. Regarding Bi-CG

Note, the standard implementation of Bi-CG does not require the storage of $V_k$ or the solution to a tridiagonal linear system; see Algorithm 2. Instead, it is based on an implicitly formed LU decomposition of $T_k$, which often has better numerical behavior compared to the computation above [19].

Similar to the method GMRES, the iterates and corresponding residual vectors are elements of a Krylov subspace, formalized as follows.

**Proposition 2.27** (Residual of iterates, Bi-CG). The $k$th iterate of Bi-CG defined in (2.12) is an element of the Krylov subspace $K_k(A,b)$. Moreover, the corresponding residual vector $r_k := b - Ax_k$ is included in $K_{k+1}(A,b)$.

The following characterizes the residual.

**Definition 2.28** (Petrov-Galerkin condition, Bi-CG). The property

$$b - Ax_k \perp K_k(A^T,c)$$

is referred to as the Petrov-Galerkin condition for the Bi-CG method.

Applying the Bi-CG method to (1.1) requires a Lanczos biorthogonalization procedure. Here the Krylov subspace basis vectors $v_{k+1}$ and $w_{k+1}$ can be computed using a linear combination of only three vectors, shown as follows.

**Corollary 2.29.** Let $v_1, \ldots, v_k$ and $w_1, \ldots, w_k$ denote the columns of $V_k$ and $W_k$, respectively. The Krylov vector $v_{k+1}$ in (2.8) fulfills a three-term recurrence between itself and the previous Krylov vectors since

$$A \begin{bmatrix} v_1 & \cdots & v_k \end{bmatrix} = \begin{bmatrix} v_1 & \cdots & v_{k+1} \end{bmatrix} \begin{bmatrix} \alpha_1 & \gamma_1 & & & \\ \beta_1 & \ddots & \ddots & & \\ & \ddots & \ddots & \ddots & \\ \gamma_{k-1} & \beta_{k-1} & \alpha_k & & \\ \gamma_k & \beta_k & & & \end{bmatrix}$$

and $Av_k = \gamma_{k-1}v_{k-1} + \alpha_kv_k + \beta_kv_{k+1}$. Similarly,

$$A^T \begin{bmatrix} w_1 & \cdots & w_k \end{bmatrix} = \begin{bmatrix} w_1 & \cdots & w_{k+1} \end{bmatrix} \begin{bmatrix} \alpha_1 & \beta_1 & & & \\ \gamma_1 & \ddots & \ddots & & \\ & \ddots & \ddots & \ddots & \\ \beta_{k-1} & \gamma_{k-1} & \alpha_k & & \\ & \gamma_k & & & \end{bmatrix}$$

and $A^T w_k = \beta_{k-1}w_{k-1} + \alpha_kw_k + \gamma_kw_{k+1}$.

Breakdown can occur in the context of Bi-CG, and the situation is not always positive. We clarify this with the following remark.
Remark 2.30 (Breakdown, Bi-CG). If $v_i = 0$ or $w_i = 0$, an invariant subspace of $A$ or $A^T$, respectively, has been found, and the algorithm terminates. If an invariant subspace with $A$ has been found, then the approximate solution to (1.1) is exact. In the case that $w_i = 0$, the approximation to the dual system $A^T z = c$ is exact, though the approximation to (1.1) may not be. Either of these situations is referred to as a lucky breakdown. A so-called serious breakdown occurs if $w_i^T v_i = 0$, for $v_i \neq 0$, $w_i \neq 0$. No further progress can be made in this scenario, and the approximate solutions are not exact.
The main scientific contribution of this thesis is the collection of four papers included in Part II. Below are short summaries of each of the appended papers, where the author’s contributions are highlighted.

### 3.1 Infinite GMRES

Paper A of this thesis is


**Summary:** Paper A proposes a Krylov subspace method to solve parameterized linear systems of the form $A(\mu)x(\mu) = b$ for many different $\mu \in \mathbb{C}\backslash\{0\}$. In particular, the matrix $A(\mu) \in \mathbb{C}^{n \times n}$ is of the form

$$A(\mu) = A_0 + A_1f_1(\mu) + \ldots + A_{nf}f_{nf}(\mu), \quad \text{for} \quad nf \ll n,$$

where the functions $f_i$ are nonlinear in $\mu$, the matrices $A_i$ are large and sparse, and the vector $b \in \mathbb{C}^n$. The nonsingular matrix $A(\mu)$ is analytic in $\mu$ and, thus, can be described by an infinite Taylor series expansion. This work considers such an expansion centered at $\mu = 0$, i.e.,

$$A(\mu) = \sum_{i=0}^{\infty} \frac{1}{i!} A^{(i)}(0) \mu^i, \quad (3.1)$$

where we assume that many of the derivatives $A^{(i)}(0)$ are computationally available and $A(0)$ is nonsingular.
Chapter 3. Summary of Results

The approach of Paper A is summarized as follows. Let $A_N(\mu) \in \mathbb{C}^{n \times n}$ denote a truncated Taylor series expansion of $A(\mu)$, consisting of only the first $N + 1$ terms in (3.1). A linearization of dimension $(N + 1)n \times (N + 1)n$ and the form

$$(\mu K_N - I_N)v(\mu) = c$$

is considered, where the matrix $K_N$ is constructed using a linear transformation of the Taylor series coefficients of $A_N(\mu)$. Furthermore, the matrix $K_N$ and the vector $c$ are independent of the parameter $\mu$, and $I_N$ denotes the identity matrix of the appropriate dimension. Solutions to (3.2) are unique for each $\mu$ and equivalent in a certain sense to solutions of the corresponding linear systems $A_Nx_N(\mu) = b$, where $A_N(\mu)$ is assumed nonsingular. Specifically, the solutions to the linearized system (3.2) are given by

$$v(\mu) := \left[\begin{array}{c} \mu^0T x_N(\mu) \\ \mu^1T x_N(\mu) \\ \vdots \\ \mu^N T x_N(\mu) \end{array}\right]^T \in \mathbb{C}^{(N+1)n},$$

proved in this work.

Applying GMRES to (3.2) requires the construction of an orthonormal basis for the Krylov subspace

$$\mathcal{K}_k(K_N, c) := \text{span} \{c, K_N c, \ldots, K_N^{k-1} c\},$$

since Krylov subspaces are shift- and scaling-invariant. The associated basis matrix is built once and reused to compute approximations to the linearized systems for each $\mu$ of interest. Moreover, the basis matrix for (3.3) is constructed using the infinite Arnoldi algorithm [31]. We show that this allows us to execute the algorithm in a finite number of operations without truncation in the Taylor series expansion, i.e., $N \to \infty$. In this way, approximating the solution to (3.2) is equivalent to approximating (1.6) with $\ell = 1$ and complex coefficients. The method is referred to as infinite GMRES.

Paper A also provides a tensor variant of infinite GMRES, inspired by [30]. This version of the algorithm offers an effective way of handling the basis matrix, which is often large. Additionally, a low-rank variant, based on [61], is developed. The low-rank variant is more computationally efficient for problems in which the higher-order terms in the Taylor series expansion of $A(\mu)$ have a reduced rank.

A linear solve with $A(0)$ and a unique right-hand side vector is required at each iteration of the infinite Arnoldi algorithm in this context, and this is performed via a precomputed LU decomposition. The convergence analysis is based on the eigenvalues of $K_N$ and the magnitude of the parameter $\mu$. Numerical simulations involving a parameterized Helmholtz equation and a transfer function of a time-delay system show the suitability of these algorithms for solving large-scale problems.

**Contribution:** The ideas in this work were suggested by Elias Jarlebring. The convergence theory presented in Section 4 was worked out in close collaboration between the two authors. The author of this thesis generated all numerical results, and she wrote the manuscript with input from the coauthor.
3.2 Preconditioned infinite GMRES

Paper B of this thesis is


Summary: Paper B proposes a novel method to solve (1.6) with $\ell = 1$, where $A(\mu) \in \mathbb{R}^{n \times n}$ is nonsingular and analytic and $b \in \mathbb{R}^n$. We assume the matrix $A(\mu)$ is of the form

$$A(\mu) = A_0 + A_1 f_1(\mu) + \ldots + A_{n_f} f_{n_f}(\mu), \quad \text{for} \quad n_f \ll n,$$

where $f_i$ are nonlinear in $\mu \in \mathbb{R}$ and the matrices $A_i$ are large and sparse. In particular, the proposed method returns approximate solutions to $x(\mu)$ for many different values of $\mu$.

Let $A_N(\mu) \in \mathbb{R}^{n \times n}$ denote the Taylor series expansion of $A(\mu)$, truncated after the first $N + 1$ terms. This work is based on a linearization of $A_N(\mu)$. Specifically, the linearization is of the form

$$(K_N - \mu M_N)\hat{v}(\mu) = c,$$

has dimension $(N + 1)n \times (N + 1)n$, and includes the Taylor series coefficients of $A_N(\mu)$. Notably, the matrices $K_N$ and $M_N$, as well as the vector $c$, are independent of the parameter $\mu$. Solutions to the nonsingular system $A_N x_N(\mu) = b$ and the linearized system (3.5) are unique for each value of $\mu$ and equivalent in the sense that

$$\hat{v}(\mu) = [\mu^0 x_N(\mu)^T \quad \mu^1 x_N(\mu)^T \quad \ldots \quad \mu^N x_N(\mu)^T]^T \in \mathbb{R}^{(N+1)n},$$

which is shown in this work. Specifically, we assume that the derivatives $A^{(i)}(0)$ are computationally available.

The linear system (3.5) is solved in an inexact GMRES setting with right preconditioning. In particular, preconditioning with the matrix $K_N$ leads to the shifted parameterized linear system

$$(I_N - \mu M_N K_N^{-1})\hat{\theta}(\mu) = c,$$

where $\hat{\theta}(\mu) = K_N \hat{v}(\mu)$ and $I_N$ is the identity matrix of the appropriate dimension. Exploiting the shift- and scaling-invariance properties of Krylov subspaces allows for the computation of solutions to (3.6) for a variety of $\mu$ simultaneously. More precisely, two basis matrices are constructed with the infinite Arnoldi algorithm [31], i.e., $N \to \infty$. Approximate solutions to (1.6) are computed from one basis matrix, and this is done in a finite number of operations.

At each infinite Arnoldi iteration, $K_N^{-1}$ must be applied to a unique vector. This results in the need to perform many different linear solves with the nonsingular system matrix $A(0)$. Specifically, iterative methods are considered, where the accuracy is relaxed as the residual of the outer method decreases. Notably, in some
simulations, the error in the application of the preconditioner can be relatively high after just a few outer iterations, without degrading the overall convergence.

Theoretical analysis leads to a bound on the residual at each iteration. The competitiveness of the approach is illustrated through numerical simulations of solutions to a parameterized Helmholtz equation, as well as evaluations of a transfer function of a time-delay system. We refer to our algorithm as inexact infinite GMRES.

**Contribution:** Exploring the use of preconditioning to solve parameterized linear systems was suggested by Elias Jarlebring. The author of this thesis derived the main algorithm independently. She designed and produced the numerical experiments with support from the coauthors, and she wrote the manuscript with input from the coauthors.

### 3.3 Preconditioned Chebyshev BiCG

Paper C of this thesis is


**Summary:** Paper C proposes two variants of a three-term recurrence Krylov subspace method to approximate solutions to (1.6) with $\ell = 1$. In particular, solutions corresponding to many different values of the parameter $\mu \in [-a, a]$, for $a \in \mathbb{R}_+$, are returned simultaneously, and it is assumed that the matrix $A(\mu)$ is of the form described in (3.4). The approach used in Paper C is summarized as follows.

A linearization based on a Chebyshev approximation of $A(\mu)$ is constructed. In particular, we consider $P(\mu) \approx A(\mu)$ with

$$P(\mu) = P_0 \tau_0(\mu) + \ldots + P_N \tau_N(\mu),$$

where $P_i \in \mathbb{R}^{n \times n}$ are coefficient matrices, $\tau_i(\mu)$ are the recursively defined Chebyshev polynomials on the interval $[-a, a]$, and $N$ is the truncation parameter. It is assumed that $N$ is chosen large enough that the error introduced by the Chebyshev interpolation is small. The linearization is of the form

$$(K_N - \mu M_N)v(\mu) = c$$

and has dimension $Nn \times Nn$. The matrices $K_N$ and $M_N$ contain the Chebyshev coefficients and $K_N$, $M_N$, and $c$ are independent of the parameter $\mu$.

This work makes use of *shift-and-invert preconditioning* with preconditioner $(K_N - \sigma M_N)$, for $\sigma \in (-a, a)$, selected because

$$(K_N - \mu M_N)(K_N - \sigma M_N)^{-1} \approx I_N$$

when $\sigma \approx \mu$. Consequently, the corresponding right preconditioned system is given by

$$(K_N - \mu M_N)(K_N - \sigma M_N)^{-1} \hat{v}(\mu) = c,$$
which can be expressed equivalently as

\[(K_N - \mu M_N + \sigma M_N - \sigma M_N)(K_N - \sigma M_N)^{-1} \hat{v}(\mu) = c \]

\[\iff (I_N + (-\mu + \sigma)M_N(K_N - \sigma M_N)^{-1}) \hat{v}(\mu) = c.\]

Here \(\hat{v}(\mu) = (K_N - \sigma M_N)v(\mu)\) and \(I_N\) is the identity matrix of dimension \(Nn \times Nn\).

The short-term recurrence method known as *preconditioned Bi-CG* is applied to the linearized system above, and the shift- and scaling-invariance properties of Krylov subspaces are exploited. Concretely, one basis matrix for

\[\mathcal{K}_k(M_N(K_N - \sigma M_N)^{-1}, c)\]

is generated and used to compute approximations to (1.6) for many different \(\mu\). This is done using the procedure described in [1, 21], where, notably, the basis matrix does not need to be stored. In this setting, the preconditioner, as well as its adjoint, are applied efficiently based on the prior works [3, 33]. We refer to this method as *preconditioned Chebyshev Bi-CG*.

A variant of preconditioned Chebyshev Bi-CG which considers an inexact application of the preconditioner \((K_N - \sigma M_N)\) is derived as well. Further, a bound on the residual of the inexact method at each iteration is shown. Numerical examples of the solution to a parameterized Helmholtz equation and the evaluation of a transfer function of a time-delay system show the competitiveness of the approach in solving large-scale systems for many values of the parameter.

**Contribution:** The idea of investigating short recurrence methods for solving parameterized linear systems was suggested by Elias Jarlebring. The author of this thesis derived the main algorithm independently, and she designed and produced all numerical experiments. The work was significantly improved from discussions with and feedback from Daniel B. Szyld. The author of this thesis wrote the manuscript with input from the coauthors.

### 3.4 Chebyshev HOPGD

Paper D of this thesis is


**Summary:** The method proposed in Paper D constructs a reduced order model to approximate parameterized linear systems of the form (1.6) with \(\ell = 2\) and a nonlinear dependence on the parameters. In this way, the model can be evaluated to approximate the solution \(x(\mu)\) for many different values of \(\mu\), where

\[
\mu = \begin{bmatrix} \mu_1 & \mu_2 \end{bmatrix}^T \in \mathbb{R}^2 \quad \text{and} \quad \mu_1 \in [a_1, b_1], \quad \mu_2 \in [a_2, b_2].
\]

The main idea of this work is summarized as follows.
Chapter 3. Summary of Results

\[
X(\mu_1, \mu_2) \in \mathbb{R}^{n \times 5 \times 5}
\]

\[
X^k(\mu_1, \mu_2) \approx \Phi_1 f_1 + \Phi_2 f_2 + \cdots
\]

\[
\text{If the approximation fails to converge adequately on the tensor corresponding to }
x(\mu_1^{11}, \mu_2^{11}), \ldots, x(\mu_1^{N}, \mu_2^{N}),
\]

Chapter 3. Summary of Results

\[
X(\mu_1, \mu_2) \in \mathbb{R}^{n \times 5 \times 5}
\]

\[
X^k(\mu_1, \mu_2) \in \mathbb{R}^{n \times 5 \times 5}
\]

\[
\approx \Phi_1^n f_1^n + \Phi_2^n f_2^n + \cdots
\]

\text{Fig. 3.1. Tensor matrix } X(\mu_1, \mu_2) \text{ with } N = 9 \text{ snapshots (dots connected by vertical lines) and approximation } X^k(\mu_1, \mu_2).

The method requires precomputed solutions } x(\mu_1, \mu_2) \text{, referred to as snapshots, corresponding to a set of nodes } \mu, \text{ where }

\[
\mu := \{(\mu_1^{11}, \mu_2^{11}), \ldots, (\mu_1^{N}, \mu_2^{N})\}.
\]

The nodes are chosen in a structured way on a sparse grid, and the snapshots \( x(\mu_1^{11}, \mu_2^{11}) \) are generated using a modified version of the method proposed in Paper C. This process is referred to as sampling, and the snapshots are assembled in a sparse tensor matrix \( X(\mu_1, \mu_2) \in \mathbb{R}^{n \times m_1 \times m_2} \), for \( m_1, m_2 \in \mathbb{N} \). Specifically, \( X(\mu_1, \mu_2) = x(\mu_1, \mu_2) \), for \( i = 1, \ldots, N \), and zero otherwise.

The reduced order model is constructed via a decomposition of \( X(\mu_1, \mu_2) \) as proposed in [37], and the resulting approximation \( X^k(\mu_1, \mu_2) \) is a sum of \( k \) rank-one tensors; see Figure 3.1. An evaluation of this model, denoted by \( x^k(\mu_1, \mu_2) \in \mathbb{R}^n \), is given by

\[
x^k(\mu_1, \mu_2) := \sum_{j=1}^{k} \Phi^n_j F^n_j(\mu_1) F^n_j(\mu_2)
\]

\[
= x^{k-1}(\mu_1, \mu_2) + \Phi^n_k F^n_k(\mu_1) F^n_k(\mu_2),
\]

where \( \Phi^n_j \in \mathbb{R}^n \), \( F^n_j \) and \( F^n_j \) are scalar functions, and \( (\mu_1, \mu_2) \in \mu \). The \( k \) terms in the summation above, referred to as modes, are obtained by an alternating directions, greedy algorithm, where the cost of each step scales linearly with the number of unknowns \( n \). Modes are added to the approximation until

\[
\frac{\|x(\mu_1, \mu_2) - x^k(\mu_1, \mu_2)\|}{\|x(\mu_1, \mu_2)\|} < \varepsilon \quad \text{for all } (\mu_1, \mu_2) \in \mu
\]

and a predetermined tolerance \( \varepsilon > 0 \). Note, \( x^k(\mu_1, \mu_2) \) is separable in the two parameters, and this model can only be evaluated for \( (\mu_1, \mu_2) \in \mu \).

In general, we cannot guarantee that (3.9) will be satisfied for a given set of snapshots, even when \( \varepsilon \) is chosen modestly. If the approximation fails to converge adequately on the tensor corresponding to

\[
\{x(\mu_1^{11}, \mu_2^{11}), \ldots, x(\mu_1^{N}, \mu_2^{N})\},
\]

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3.4. Chebyshev HOPGD

the approach proposed in Paper D offers a way to generate a different set of snapshots in the same parameter space. Notably, this is done with little extra computation. As generating the snapshots is the dominating part of the algorithm, and it is not known a priori if a given tensor decomposition will succeed, this is an advantageous feature.

Once the model described by (3.8) has been constructed, the one variable functions $F_1^j$, $F_2^j$ can be interpolated to approximate $x(\mu_1, \mu_2)$, for $(\mu_1, \mu_2) \notin \mu$. Evaluating the resulting model requires, essentially, a linear combination of vectors of dimension $n$. Thus, the proposed method is one way to overcome the so-called curse of dimensionality associated with solving parameterized linear systems of the form (1.6).

Numerical examples arising from both a parameterized Helmholtz equation and a parameterized advection-diffusion equation are included. Simulations show the effectiveness of this approach in approximating the solution to parameterized PDEs and solving parameter estimation problems. Moreover, this method can be generalized for solving (1.6) with $\ell = 3, 4, \ldots$. We refer to our algorithm as Chebyshev HOPGD with sparse grid sampling.

**Contribution:** The author of this thesis came up with the idea for this work. She derived the algorithm and designed and produced the numerical experiments with support from Kirk M. Soodhalter. The simulations section was extended as a result of discussions with Melina A. Freitag. The author of this thesis wrote the manuscript with input from the coauthors.
References


REFERENCES


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They might be bigger, but we’re faster and never scared

Taylor Swift
Part II: Scientific Papers